A Stretching Algorithm for Parallel Real-time DAG Tasks on Multiprocessor Systems
Manar Qamhieh, Laurent George, Serge Midonnet

To cite this version:

HAL Id: hal-01090627
https://hal-upec-upem.archives-ouvertes.fr/hal-01090627
Submitted on 3 Dec 2014

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers. L’archive ouverte pluridisciplinaire HAL, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d’enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.
A Stretching Algorithm for Parallel Real-time DAG Tasks on Multiprocessor Systems

Manar Qamhieh
Université Paris-Est
LIGM
manar.qamhieh@univ-paris-est.fr

Laurent George
Université Paris-Est
LIGM/ESIEE
laurent.george@univ-mlv.fr

Serge Midonnet
Université Paris-Est
LIGM
serge.midonnet@univ-paris-est.fr

ABSTRACT

Parallelism is becoming more important nowadays due to the increasing use of multiprocessor systems. In this paper, we study the problem of scheduling periodic parallel real-time Directed Acyclic graph (DAG) tasks on $m$ homogeneous multiprocessor systems. A DAG task is an example of inter-subtask parallelism. It consists of a collection of dependent subtasks under precedence constraints. The dependencies between subtasks make scheduling process more challenging. We propose a stretching algorithm applied on each DAG tasks to transform them into a set of independent sequential threads with intermediate offsets and deadlines. The threads obtained with the transformation are two types, (i) fully-stretched master threads with utilization equal to 1 and (ii) constrained-deadline independent threads. The fully-stretched master threads are assigned to dedicated processors and the remaining processors $m' \leq m$, are scheduled using global EDF scheduling algorithm. Then, we prove that preemptive global EDF scheduling of stretched threads has a resource augmentation bound equal to $\frac{2+\sqrt{5}}{3}$ for all tasksets with $n < \phi \cdot m'$, where $n$ is the number of tasks in the taskset and $\phi$ is the golden ratio.

Keywords

Real-time scheduling, parallel tasks, Directed Acyclic Graphs, global EDF scheduling, speedup factor

1. INTRODUCTION

Increasing the performance of execution platforms has been done by increasing the speed of uniprocessor systems associated to the reduction of the size of chips leading to heating problems. Multiprocessor systems have been seen as one solution to overcome these physical limitations, by increasing execution capabilities with processor parallelism. Many practical examples of shifting towards multiprocessors can be found, such as the Intel Xeon, ARM and Cavium processor families.

However, current real-time software APIs are not able to efficiently take advantage of multiprocessor platforms especially when parallelism is considered. In the industry, the majority of designed applications are targeting uniprocessor systems. But this is expected to change in the coming few years which will focus on parallel programming to take advantage of multiprocessor architectures. There are many models of parallel application, but in this work, we are interested in a particular family of parallelism called inter-subtask parallelism, in which a parallel task consists of a collection of subtasks under precedence constraints and dependencies between subtasks. There are many parallel task models based on this family, such as the Fork-join model which is the base of the famous parallel programming API OpenMP, and the multi-threaded Segment model. But the most general model is the Directed Acyclic Graph (DAG) model, which we consider as our task model in this paper.

Real-time scheduling of DAG tasks in particular and parallel tasks in general is a challenging problem. In hard real-time systems, the correctness of a system does not only depend on the correctness of the results, but on the respect of tasks timing parameters. Real-time systems on both uniprocessor and multiprocessor systems have been studied in the last decade, and many researches and scheduling algorithms have been proposed for such platforms [6]. However, the extension of real-time scheduling w.r.t. parallel tasks with dependencies is not trivial. The scheduling analysis of such systems can be divided into two main categories: either by transforming task precedence constraints into independent sequential tasks that execute on multiprocessor systems, or by scheduling parallel tasks directly using adapted scheduling algorithms. The first method simplifies scheduling at the price of losing some of the characteristics of parallel tasks, as it removes subtasks dependencies such that classical scheduling algorithms can then be used. In this work, we study the scheduling of DAG tasks using DAG transformation with a stretching approach.

It is worth mentioning here that this work is a generalization of the stretching algorithm proposed by Lakshmanan et. al. in [9]. Their work targeted Fork-join parallel tasks and it aims at scheduling tasks as sequentially as possible by stretching their master threads up to their deadline. They proposed to schedule stretched tasks using partitioned FBB-

\footnote{A master thread of a parallel task is defined as the sequential path in the task with the longest execution time.}
FFD DM-decreasing scheduling algorithm. In our work, we consider a more general task model which is the DAG task and we provide a resource augmentation bound for global Earliest Deadline First (EDF) scheduling of stretched tasks.

The remainder of this paper is organized as follows. Section 2 presents related works w.r.t. the problem of scheduling real-time tasks on multiprocessor systems especially for parallel DAG model. Then we present in Section 3 the considered task model. In Sections 4 and 5, we explain the DAG stretching algorithm in detail and we provide a resource augmentation bound (speedup factor) analysis. Simulation results are provided in Section 6 to study the performance of our stretching algorithm. Finally, Section 7 concludes this work.

2. RELATED WORK

The scheduling of parallel real-time tasks of different models has been studied on both uniprocessor and multiprocessor systems. In the case of uniprocessor systems, a classical approach in the state-of-the-art is to transform a parallel task into a chain and to assign each subtask of the chain extra timing parameters used for their scheduling. For example, in [4], the authors considered a hybrid taskset of periodic independent tasks and sporadic dependent graph tasks with a DAG model. They proposed an algorithm aiming at modifying the DAG timing parameters (by adding intermediate offsets and deadlines) in order to remove the dependencies between the tasks in the analysis.

In the case of multiprocessor systems, most research has been done regarding scheduling hard real-time tasks on homogeneous multiprocessor systems, especially for independent sequential tasks [6]. As mentioned above, there are mainly two methods for scheduling parallel DAGs in hard real-time systems. The first method schedules DAG tasks by using directly common scheduling algorithms and adapting the performance analysis and the scheduling conditions to take into consideration the particular characteristics of DAGs and parallel tasks in general. This technique is introduced in [2], which considers a taskset of a single sporadic DAG. They also provided polynomial and pseudo-polynomial schedulability tests for EDF scheduling algorithms.

The problem of scheduling multiple DAG tasks on multiprocessors have been more studied in [3, 10]. The authors considered general timing parameters of DAG tasks without taking into account their internal structure. They proved that global EDF has a resource augmentation bound of $2 - 1/m$, where $m$ is the number of processors in the system. In [11], the internal structure and the dependencies between subtasks are considered in the analysis of global EDF scheduling of DAG tasks.

The second method for DAG scheduling is based on DAG transformation. Dependencies of inter-subtask parallelism are removed and a DAG task is transformed into a collection of independent sequential threads. A decomposition algorithm is provided in [12] to distribute the slack time of the DAG task on its subtasks. The slack time is defined as the difference between the deadline of the graph and its minimum sequential execution time. The subtasks are assigned intermediate offsets and deadlines. The authors proved that preemptive global EDF scheduling algorithm has a resource augmentation bound equal to 4, and 4 plus a constant non-preemption overhead for non-preemptive global EDF.

In [9], the authors considered fork-join model of parallel tasks and they proposed a stretching algorithm to execute them as sequentially as possible. A fork-join model is represented as an alternative sequence of sequential and parallel segment. All parallel segments of the same task have the same number of threads, and the threads of each segment have the same sequential execution length. The authors proposed to stretch a fork-join task into a master thread with utilization equal to 1, the remaining parallel threads are forced to execute in parallel with the master thread within a fixed activation interval. Hence, dependencies are no longer considered in the scheduling process. A resource augmentation bound of 3.42 for partitioned preemptive FBB-FFD using Deadline Monotonic (DM) scheduling is given.

Fork-join model is a special case of DAG task model, our paper is an extension of the stretching algorithm of fork-join tasks to DAG tasks. The concept of our stretching algorithm is the same as the one proposed in [9], but the steps and the analysis are adapted to a general DAG task model. The contributions of this paper are as follows:

- we propose a stretching algorithm adapted to DAG task model. This algorithm is a pre-step of the scheduling process of the DAGs, which transforms parallel DAGs into independent sequential threads to be scheduled on $m$ identical processors;
- for a taskset of $n$ DAG tasks that execute on a platform of $m$ identical processors, we prove that global EDF scheduling of stretched tasks has a resource augmentation bound equal to $\frac{3+\sqrt{5}}{2}$ for all tasks if $n < \varphi \cdot m'$, where $m' \leq m$ is the number of processors not running master threads and $\varphi$ is the golden ratio, and equal to 4 else.

3. TASK MODEL

We consider a taskset $\tau$ of $n$ periodic parallel real-time Directed Acyclic Graph (DAG) tasks run on a system of $m$ identical processors. The taskset $\tau$ is represented by $\{\tau_1, ..., \tau_n\}$. Each DAG task $\tau_i$, where $1 \leq i \leq n$, is a periodic implicit-deadline graph which consists of a set of subtasks under precedence constraints that determine their execution flow. A DAG task $\tau_i$ is characterized by $\{\{\tau_{i,j} | 1 \leq j \leq n_i\}, G_i, O_i, D_i\}$, where the first parameter represents the set of subtasks of $\tau_i$, and $n_i$ is their number. $G_i$ is the set of directed relations between these subtasks, $O_i$ is the offset of the DAG and $D_i$ is $\tau_i$'s relative deadline. Since each DAG task has an implicit deadline, its period $T_i$ (interval time between its successive jobs) is the same as its deadline $T_i = D_i$.

Let $\tau_{i,j}$ denotes the $j^{th}$ subtask of the set of subtasks forming the DAG task $\tau_i$, where $1 \leq j \leq n_i$. Each subtask $\tau_{i,j}$ is a single-threaded sequential task which is characterized by a worst-case execution time (WCET) $C_{i,j}$. All subtasks of a DAG share the same deadline and period of the DAG. The total WCET $C_i$ of DAG $\tau_i$ is defined as the sum of WCETs...
Figure 1: An example of a DAG task $\tau_1$ which consists of 7 subtasks. The number on the upper right corner of each subtask represents its WCET and the arrows represent their precedence constraints.

of its subtasks, where $C_i = \sum_{j=1}^{n_i} C_{i,j}$. Let $U_i$ denote the utilization of $\tau_i$, where $U_i = \frac{C_i}{T_i}$, and its density is denoted by $\delta_i = \frac{C_i}{\min(D_i,T_i)}$. For an implicit-deadline task, whose deadline is equal to its period, the density is the same as the utilization, while the density of a constrained-deadline task becomes $\delta_i = \frac{C_i}{D_i}$.

The set of directed relations between the subtasks of DAG $\tau_1$ is denoted by $G_1$ and it defines their dependencies. A directed relation $G_i(j,k)$ between subtasks $\tau_{i,j}$ and $\tau_{i,k}$ means that $\tau_{i,j}$ is a predecessor of $\tau_{i,k}$, and the latter subtask have to wait for all of its predecessors to complete their execution before it can start its own. An example of a DAG task is shown in Figure 1, in which $\tau_1$ consists of 7 subtasks. Precedence constraints are represented by directed arrows between subtasks. A source subtask is a subtask with no predecessors such as $\tau_{1,1}$. Respectively, an ending subtask is the one without any successors such as $\tau_{1,7}$.

Based on the structure of DAG tasks, the critical path of DAG $\tau_1$ is defined as the longest sequential execution path in the DAG when it executes on a virtual platform composed of an infinite number of processors. Its length $L_i$ is the minimum response time of the DAG. A subtask that is part of the critical path is referred to as a critical subtask, while non-critical subtasks are executed in parallel with the critical ones.

A DAG task is said to be feasible if the subtasks of all of its jobs respect its deadline. A taskset $\tau$ is deemed unfeasible when scheduled using any scheduling algorithm on $m$ unit-speed processors if, at least, one of the following conditions is false:

$$\forall \tau_i \in \tau, L_i \leq D_i$$

$$U(\tau) = \sum_{i=1}^{n} U_i = \sum_{i=1}^{n} \frac{C_i}{T_i} \leq m$$

4. DAG STRETCHING ALGORITHM

Based on our task model, subtasks of a DAG inherit the absolute deadline of the DAG task. However, only source subtasks are activated by the activation of the DAG task, while their successors are activated dynamically based on the completion time of their predecessors. During scheduling process, the subtasks of a DAG have to execute within the activation interval of the DAG (between the release time and the absolute deadline of the DAG), while the exact activation time of subtasks is unknown prior to scheduling process. As a result, multiprocessor scheduling of DAG tasks is more complicated than the scheduling of independent sequential tasks. Therefore, scheduling process of DAG tasks can be simplified by avoiding the dynamic activation of subtasks by assigning them intermediate offsets and deadlines.

Hereby, we propose a stretching algorithm for parallel DAG tasks which transforms each DAG task into a set of independent periodic constrained-deadline threads, that can be scheduled using any scheduling algorithm of multiprocessors. With our DAG stretching algorithm, the scheduling of a DAG task is done based on the timing parameters of the threads, and a task is deemed feasible if all of its jobs respect their assigned intermediate deadlines. Our DAG stretching algorithm chooses intermediate offset and deadline for each subtask in a DAG task. Before explaining the concept of our algorithm, we start by analyzing the DAG model and identifying its characteristics that led to the stretching algorithm.

Assume that the critical subtasks of DAG $\tau_1$ are combined together in a sequential master thread $\tau_1^{master}$ whose execution time is equal to $L_1$. Based on the task model, $\tau_1$ can be seen as a master thread with the longest sequential execution time among all paths in the DAG from a source to an ending subtask, while the remaining paths of the DAG execute in parallel with it. For example, DAG task $\tau_1$ from Figure 1 consists of 7 subtasks and 6 execution paths: $\{\{\tau_{1,1}, \tau_{1,4}, \tau_{1,6}\} \{\tau_{1,2}, \tau_{1,4}, \tau_{1,6}\} \{\tau_{1,3}, \tau_{1,4}, \tau_{1,6}\} \{\tau_{1,5}, \tau_{1,4}, \tau_{1,6}\}\}$.

The slack time of the DAG can be seen as a slack time of a master thread $\tau_1^{master}$ when it is scheduled exclusively without interference from other tasks. $Sl_i$ denote the positive slack available to DAG $\tau_1$ when it is scheduled exclusively without interference from other tasks. $Sl_i$ is given by:

$$Sl_i = D_i - L_i$$

The slack time of the DAG can be seen as a slack time of its master thread since it is the DAG’s longest path. Back to the previous example, if we assume that $\tau_1$ had a deadline $D_1 = 10$, then its slack time would be $Sl_1 = 4$. In order to avoid dependencies between the subtasks of a DAG task, we propose to use the stretching algorithm which is summarized as follows:

**DAG Stretching Algorithm.**

Our stretching algorithm fills the slack time of a DAG task by adding fractions of non-critical subtasks to its master thread until it is stretched up to DAG’s deadline. As a result, the remaining non-critical subtasks (if any) will be forced to execute in parallel with the stretched master thread within a fixed activation interval.

The objective of our stretching algorithm is to avoid the
parallel structure of DAG tasks by executing their subtasks as sequentially as possible. Hence, for each DAG task \( \tau_i \in \tau \), if \( \tau_i \) fits completely on a single processor \((C_i \leq D_i)\), then the stretching algorithm transforms it into a single thread which contains all of its subtasks and forces them to execute sequentially. Otherwise, the algorithm fully stretches the master thread of \( \tau_i \) up to its deadline \( D_i \). As a result, the stretching algorithm generates a fully-stretched master thread and a collection of independent threads with intermediate offsets and deadlines that execute in parallel with the master thread. The offsets and deadlines of threads are important for the scheduling process and also to maintain the precedence constraints of DAG \( \tau_i \). It is worth noting that the stretching algorithm is a pre-step to the scheduling process, and any master thread with utilization equal to 1, is assigned a dedicated processor by the scheduler. Other parallel threads are scheduled on the remaining processors of the system using any multiprocessor scheduling algorithm. In this paper, we choose global EDF scheduling algorithm.

For clarity reasons, the reader is advised to refer to the example in Figures 2 and 3 so as to have a better understanding of our DAG stretching algorithm. More details about the example are provided at the end this section.

In order to perform our stretching algorithm, we propose to transform a DAG task into Multi-Threaded Segment (MTS) representation. This is a basic transformation for DAGs that maintains precedence constraints between their subtasks and leads to an easier scheduling analysis.

**The Multi-Threaded Segment (MTS) Representation**

A multi-threaded Segment (MTS) task \( \tau'_i \) consists of a sequence of segments. Each segment is composed of a number of parallel threads with the same WCET. Let \( S_i \) denote the set of parallel segments of \( \tau'_i \), and \( s_i \) denote the total number of segments in \( S_i \). Each segment \( S_{i,j}, 1 \leq j \leq s_i \), consists of \( m_{i,j} \) parallel threads. All threads of segment \( S_{i,j} \) have the same WCET which is denoted by \( e_{i,j} \). The number of threads and their WCETs vary from one segment to another, and each segment contains at least one thread.

In this work, MTS task \( \tau'_i \) is not a new task but a representation of DAG task \( \tau_i \in \tau \). This is done by considering that all subtasks of \( \tau_i \) are executed as soon as possible on a virtual platform having an infinite number of processors. As a result, each subtask executes under precedence constraints and it does not suffer from interference of other subtasks in the system. For a DAG task \( \tau_i \), its source subtasks are activated at the activation time of DAG \( \tau_i \). Then, successor subtasks are activated as soon as their predecessors have completed their execution. A segment in the MTS task is defined whenever a subtask completes its execution. Thus, threads of the same segment have the same WCET. It is worth noticing that a subtask in DAG \( \tau_i \) may be divided into two or more threads executing in successive segments from its MTS representation \( \tau'_i \).

Figure 2 shows the MTS representation \( \tau'_i \) of DAG \( \tau_i \) from Figure 1. Knowing that DAG \( \tau_i \) is released at time \( t = 0 \) and has a deadline \( D_i = 10 \), source subtasks \{\( \tau_{i,1}, \tau_{i,2}, \tau_{i,3}, \tau_{i,5} \}\) are activated at \( t = 0 \). At \( t = 2 \), both subtasks \( \tau_{i,3} \) and \( \tau_{i,5} \) complete their execution and the first segment \( S_{i,1} \) is defined. Their successors (subtasks \( \tau_{i,6} \) & \( \tau_{i,7} \)) have to wait for all of their predecessors to complete their execution before they can start theirs. At time \( t = 3 \), subtasks \( \tau_{i,1} \) and \( \tau_{i,2} \) complete their execution, subtask \( \tau_{i,4} \) starts its own and segment \( S_{i,2} \) is defined. Finally, at time \( t = 4 \), subtasks \( \tau_{i,6} \) and \( \tau_{i,7} \) are activated, and segments \( S_{i,3}, S_{i,4} \) and \( S_{i,5} \) are defined at times 4, 5 and 6 respectively. The resulting MTS task \( \tau'_i \) is represented as a sequence of 5 segments. The number and the WCET of threads of each segment are identified. Subtask \( \tau_{i,1} \) is an example of a subtask that is spread on multiple segments. It is divided into two threads executing in segments \( S_{i,1} \) and \( S_{i,2} \).

Finally, the MTS task \( \tau'_i \) shares the same deadline \( D_i \) and period \( T_i \) of DAG task \( \tau_i \). Hence, Equation 1 regarding the slack \( S_i \) of \( \tau_i \) remains correct for the MTS task \( \tau'_i \). Also, the critical path length \( L_i \) and the total WCET \( C_i \) of the original DAG task \( \tau_i \) is the same for \( \tau'_i \), but they are calculated differently based on the parameters of the MTS task \( \tau'_i \):

\[
L_i = \sum_{j=1}^{s_i} e_{i,j}
\]

\[
C_i = \sum_{j=1}^{s_i} m_{i,j} \cdot e_{i,j}
\]

For a given DAG set \( \tau \), the stretching algorithm is applied to \( \tau'_i \) for each DAG task \( \tau_i \in \tau \). The algorithm is explained in detail in the remainder of this section.

**The DAG Stretching Algorithm**

As stated earlier, the DAG stretching technique executes the DAGs as sequentially as possible. By doing so, parallel structure of DAGs is removed and dependencies between the subtasks are replaced by intermediate offsets and deadlines. The stretching algorithm is done based on the MTS representation \( \tau'_i \) rather than the DAG structure \( \tau_i \). Hence, Equation 1 regarding the slack \( S_i \) of \( \tau_i \) remains correct for the MTS task \( \tau'_i \). Also, the critical path length \( L_i \) and the total WCET \( C_i \) of the original DAG task \( \tau_i \) is the same for \( \tau'_i \), but they are calculated differently based on the parameters of the MTS task \( \tau'_i \):

- if \( U_i \leq 1 \), then DAG \( \tau_i \) is considered as a sequential task because its total WCET \( C_i \) can be contained entirely within its deadline \( D_i \). As a result, the stretching algorithm transforms its subtasks into a single master thread \( \tau_{i,master} \) in which all of the subtasks execute sequentially. If \( U_i \) is equal to 1, then the master thread...
a fully-stretched master thread. In order to do so, parallel non-critical threads of \( \tau_i \) are used to fill the slack \( Sl_i \) of the task uniformly, and all segments of \( \tau_i' \) with more than one thread participate. To achieve a uniform filling of the slack, equal fractions from these threads are added to the master thread. In this case, where the stretching algorithm targets only parallel tasks with \( C_i > D_i \), the total WCET \( C_i \) of \( \tau_i' \) without its critical path length \( L_i \) definitely exceeds its deadline \( D_i \). We define a distribution factor \( f_i \) of each execution unit in \((C_i - L_i)\) that has to be added to the slack \( Sl_i \) as follows:

\[
Sl_i = \frac{D_i - L_i}{C_i - L_i} \leq \frac{D_i}{C_i} < 1
\]

In order to clarify the meaning of \( f_i \), let us suppose that each non-critical execution unit in task \( \tau_i' \) is divided into two parts, first part of length equal to \( f_i \) is added to the master thread and \((1 - f_i)\) execute in parallel as a constrained-deadline thread. The total execution time added from non-critical threads to the master thread is equal to \( Sl_i \). However, it is impractical to use \( f_i \) in this way, and force each execution unit to be divided into two parts. So, the filling of the slack is based on the execution requirement of each segment in the task and not on the threads directly. We fill the slack with the maximum number of entire threads from each segment. Thus, at most one thread from each segment will be used to fill the slack partially. From the MTS representation \( \tau_i' \), each segment \( S_{i,j} \in \tau_i \) has \( m_{i,j} \) parallel threads, among them there is a critical thread. Hence, the total non-critical worst-case execution time of segment is equal to \((e_{i,j} + (m_{i,j} - 1))\). Based on the definition of \( f_i \), we conclude that a total of \((f_i * e_{i,j} (m_{i,j} - 1))\) time units from segment \( S_{i,j} \) will be added to the master thread.

Now, we want to identify how many threads of each segment \( S_{i,j} \) are added to the master thread based on the total execution time of the segment. Knowing that the threads of segment \( S_{i,j} \) have equal WCETs denoted by \( e_{i,j} \), which is equivalent to the shortest sequential execution length of the segment, we can identify how many entire threads are added to the master thread. Let \( f_{i,j} \) denote the number of threads from segment \( S_{i,j} \) to be added to the master thread:

\[
f_{i,j} = f_i * e_{i,j} (m_{i,j} - 1) / e_{i,j} = f_i * (m_{i,j} - 1)
\]

According to this, each segment \( S_{i,j} \) adds \([f_{i,j}]\) entire threads and a fraction of a thread of length \((f_{i,j} + [f_{i,j}])e_{i,j}\) to the master thread. As a result, the slack \( Sl_i \) of task \( \tau_i' \) is filled completely and a fully-stretch master thread \( \tau_i' \) with \( U_i^{\text{master}} = 1 \) is generated. We conclude that each segment \( S_{i,j} \) adds in total \((1 + f_i)\) threads to the master thread (including the critical thread), while the remaining threads of the segment execute in parallel with the master thread. Hence, each segment \( S_{i,j} \) has an intermediate deadline \( D_{i,j} \) calculated as follows:

\[
D_{i,j} = (1 + f_{i,j}) * e_{i,j}
\]

From the definition of the MTS model, segments of a task \( \tau_i' \) execute sequentially and when one segment completes its execution, its successor starts its own. Hence, at any time
According to this, we can define an intermediate offset $O_{i,j}$ for each segment $S_{i,j} \in \tau'$ based on the intermediate deadlines of the segments, where:

$$\forall S_{i,j} : j > 1 \rightarrow O_{i,j} = \sum_{k=1}^{j-1} D_{i,k}$$

and $O_{i,1} = 0$ (since $\tau_1$ has no offset).

After applying the stretching algorithm, a segment $S_{i,j}$ of $\tau'$ comprises of:

- a thread $\tau^\text{master}_{i,j}$ which is part of the master thread $\tau^\text{master}_i$ of $\tau'$. It has a WCET of $D_{i,j}$ and a deadline of $D_{i,j}$.
- $(m_{i,j} - [f_{i,j}] - 2)$ parallel constrained-deadline threads with a WCET of $e_{i,j}$ and a deadline $D_{i,j}$.
- one remaining thread with WCET of $(1 + [f_{i,j}] - f_{i,j})e_{i,j}$ and a constrained-deadline $(1 + [f_{i,j}])e_{i,j}$. The remaining WCET is added to the master thread so as to fill its slack. We can notice here that the remaining thread has a shorter deadline than the other remaining threads from the same segment. This is done so as to force the first fraction of the thread to finish its execution before the second fraction in the master thread starts its own. As a result, both fractions of the same threads are forced to execute sequentially.

For each segment $S_{i,j}$, the total number of threads in a segment (including the partial thread), which are not added to the master thread, is given by:

$$q_{i,j} = m_{i,j} - [f_{i,j}] - 1$$

According to this, $q_{i,j}$ independent threads from segments of each task $\tau' \in \tau$ are scheduled using any multiprocessor scheduling algorithm, while the fully-stretched master threads are assigned their own processors. After applying the stretching algorithm on each task $\tau' \in \tau$, each task generates at most one fully-stretched master thread. Hence, the total number of these master threads, respectively the number of their dedicated processors, cannot exceed $n$, which is the total number of tasks in the original DAG set $\tau$. As a result, the constrained-deadline tasks are scheduled on the remaining processors of the system which are referred to as $m'$. The relation between the remaining processors $m'$ and the total number of processors in the system $m$ is given as follows:

$$m' \geq m - n$$

Algorithm 1 shows the DAG stretching algorithm. It shows its steps and the generated threads after the stretching. Now, we present an example of the stretching algorithm applied on the DAG task $\tau_1$ from Figure 1.

**Example: Stretching Algorithm**

We present in this section an example of the stretching algorithm applied to the DAG task $\tau_1$ of Figure 1. The algorithm is summarized in three main steps which are shown in Figures 2 & 3. We assume that task $\tau_1$ is a periodic implicit-deadline DAG with $C_1 = 14$ and a deadline $D_1 = 10$. Since its utilization $U_1 = 1.4 > 1$, then $\tau_1$ has to be represented by $\tau'_1$ which is shown in Figure 2. The subtasks of $\tau_1$ execute as soon as possible while considering an execution platform of infinite number of processors, and the only blocking effect on a subtask is due to its predecessors.

In Figure 2, the MTS task $\tau'_1$ is identified. It consists of 5 segments. Segment $S_{1,1}$ has 4 threads with $e_{1,1} = 2$, while the remaining threads have WCET equal to 1. Segments $S_{1,2}$ & $S_{1,4}$ have two threads and segments $S_{1,3}$ & $S_{1,5}$ have a single thread. The length of the critical path is $L_1 = 6$ and its deadline $D_1 = 10$. Based on Equation 1, the slack $S_1$ is equal to 4. The total WCET of $\tau_1$ is equal to 14 and the non-critical execution time (while excluding the critical path) is equal to 8. According to Equation 2, the distribution factor of $\tau'_1$ is equal to $f_1 = \frac{4}{8} = \frac{1}{2}$. For each segment $S_{i,j}$, its distribution factor $f_{i,j}$ depends on the number of its parallel
threads, where \( f_{1,1} = \frac{3}{2} \), \( f_{1,2} = f_{1,4} = \frac{1}{2} \) and \( f_{1,3} = f_{1,5} = 0 \), as shown in Figure 3(a). It is worth noticing that segments with a single thread have a distribution factor equal to zero, because their single thread is a critical one and it is already included in the master thread.

Based on the values of the segment distribution factor of each segment, we can identify how many threads to be added to the master thread from each segment. As described earlier, \( 1 + f_{i,j} \) threads from segment \( S_{i,j} \) are used to fill the master thread. Figure 3(b) shows the stretching of task \( \tau_1 \) and its final result. We can notice that a total of \( \frac{5}{2} \) threads from segment \( S_{1,1} \) are added to the master thread where each thread has a length of 2. This is equivalent to 2 entire threads and a half. Knowing that the original DAG task \( \tau_1 \) has no offset, then the first segment \( S_{1,1} \) has no offset and its execution interval length is defined as \( \frac{5}{2} + 2 = 5 \). The third thread of segment \( S_{1,1} \) (denoted by \( \tau_{1,3} \) in Figure 3(b)) is divided into two halves, one is added to the master thread and the other is an independent constrained-deadline thread with a deadline \( D_{1,3} = 4 \). The fourth remaining thread of the segment has a deadline equal to the deadline of the segment which is 5. The difference between both threads is that thread \( \tau_{1,3} \) has to finish earlier than the deadline of the segment so as to give its other part (the one added to the master thread) enough time to execute. Both parts have to execute sequentially and never in parallel, and the intermediate deadline of the remaining thread forces the sequential execution of the entire subtask. By applying the same calculations on all of the segments, we obtain the result in Figure 3(b). Regarding segments \( S_{1,2} \) and \( S_{1,4} \), each one adds \( \frac{1}{2} \) thread to the master thread. This means that there is only one partial thread left to execute in parallel with the master thread from each segment.

As shown in the example in Figure 3(b), the stretching algorithm generates a master thread \( \tau_{\text{master}}^{m\text{ast}} \) with WCET and deadline equal to 10, and 4 constrained-deadline threads identified as follows: \( \{ \tau_i^{cd} \} = \{ \tau_{1,3} : (0, 1, 4, 10), \tau_{1,2} : (0, 2, 5, 10), \tau_{1,5} : (5, 0, 5, 1, 10) \} \), where each thread is identified by its intermediate offset, WCET, intermediate deadline and period. These threads are scheduled as independent constrained-deadline threads on \( m' \) processors. In this example, we consider that taskset \( \tau \) contains a single DAG \( \tau_1 \) which execute on a system of 2 identical processors. Then the master thread occupies a processor for itself, and the other processor is used for the scheduling of the parallel threads, which means \( m' = 1 \).

**Lemma 1.** A DAG set \( \tau \) that is schedulable using algorithm \( A \) on a system of \( m \) processors, remains schedulable after stretching, i.e., the DAG stretching algorithm does not affect its schedulability.

**Proof.** As described above, the stretching algorithm assigns intermediate offsets and deadlines for threads of each DAG in \( \tau \), without changing their original offsets and deadlines. Intermediate offsets and deadlines maintain precedence constraints of the original DAG, which means that if a \( \tau' \) is schedulable while respecting the intermediate deadlines of the threads, then the original DAG taskset \( \tau \) must be schedulable while respecting the DAGs’ deadlines. \( \square \)

5. RESOURCE AUGMENTATION BOUND ANALYSIS

In this section, we analyze the performance of DAG stretching algorithm of parallel tasks by calculating its resource augmentation bound (speedup factor) when global EDF is used to schedule the non fully-stretched threads generated from the stretching algorithm. Speedup factor is a common performance metric for real-time scheduling algorithms, that gives indications about their performance w.r.t. the performance of an optimal algorithm.

**Definition 1 (from \([8]\)).** For a given taskset \( \tau \) that is feasible on \( m \) unit-speed processors using an optimal scheduler, it is schedulable using \( A \) scheduling algorithm on \( m' \) processors that are \( \nu \) times faster. The minimum speedup factor of processors speed is the resource augmentation bound (or simply the speedup factor) of scheduler \( A \).

Let \( \tau'' \) be the thread set generated after applying our stretching algorithm on every DAG task in \( \tau \). It contains all threads generated from the stretching algorithm except for the fully-stretched master threads (including the constrained-deadline parallel threads and the implicit-deadline master threads). We prove that global EDF scheduling of \( \tau'' \) when executing on \( m' \) processors, has a speedup factor of \( \frac{m' + \sqrt{m'} - \sqrt{m - 1}}{m' - 1} \) for all tasks with \( n < \varphi \cdot m' \), where \( \varphi = \frac{1 + \sqrt{5}}{2} \), the remaining number of processors is denoted by \( m' \geq m - n \), where \( m \) is the original number of processors in the system and \( n \) is the number of tasks in the set \( \tau \). This implies that if a taskset \( \tau'' \) is feasible on \( m' \) unit-speed processors, then it is schedulable using global EDF on \( m' \) processors with speed of \( \frac{m' + \sqrt{m'}}{m'} \) as fast as the original. Otherwise, the speedup factor is equal to 4 if we use the decomposition algorithm from \([12]\). The fully-stretched master threads are assigned their own processors, while the remaining parallel threads are scheduled using global EDF on the remaining processors.

Our speedup factor analysis is based on the following sufficient scheduling condition of global EDF scheduling algorithm.

**Theorem 1 (From \([1]\)).** Any constrained-deadline sporadic sequential task set \( \tau \) with total density \( \delta^{\text{cum}}(\tau) \) and maximum density \( \delta^{\text{max}}(\tau) \) is schedulable using preemptive global EDF policy on \( m \) unit-speed processors if

\[
\delta^{\text{cum}}(\tau) - \frac{m - (m - 1)\delta^{\text{max}}(\tau)}{m} \leq 0
\]

**Theorem 2.** If constrained-deadline taskset \( \tau'' \) of parallel threads, with \( n < \varphi \cdot m' \), where \( \varphi = \frac{1 + \sqrt{5}}{2} \) (the golden ratio) and \( n \) is the number of tasks in DAG taskset \( \tau \), is feasible on \( m' \) unit-speed processors, then \( \tau'' \) is schedulable using global EDF on \( m' \) processors with speed at least \( \frac{1 + \sqrt{5}}{2} \).

**Proof.** Based on Theorem 1, a taskset of constrained-deadline parallel threads \( \tau'' \) that is generated after applying the DAG stretching algorithm is schedulable on the remaining processors \( m' \) of the system if

\[
\delta^{\text{cum}}(\tau'') - \frac{m' - (m' - 1)\delta^{\text{max}}(\tau'')}{m'} \leq 0
\]

We start by calculating the maximum density \( \delta^{\text{max}}(\tau'') \) of the threads of \( \tau'' \). As stated earlier, \( \tau'' \) consists of two types of threads: implicit-deadline master threads (not fully-stretched) from DAG tasks with utilization less than 1, and a collection of constrained-deadline threads \( \{ \tau^{cd} \} \) from each DAG task with a utilization higher than 1. We will consider both cases while calculating \( \delta^{\text{max}}(\tau'') \):
• case 1: for an implicit-deadline DAG task $\tau_i$ with $U_i < 1$, it is completely transformed into a single master thread $\tau_i^{master}$ with the same timing parameters of the original DAG $\tau_i$. Based on the assumption that $\tau_i$ has implicit deadline, then its deadline $D_i^{master}$ is equal to its period $T_i^{master}$, and respectively, $\delta_i^{master} = U_i^{master} < 1$.

• case 2: for an implicit-deadline DAG task $\tau_i$ with $U_i > 1$, it is transformed into a set of constrained-deadline parallel threads with intermediate offsets and deadlines. From the structure of multi-threaded segment tasks, only one segment is active at every time instant $t$. Each segment $S_{i,j} \in S_i \in \tau_i, \jmath$, with $m_{i,j} > 1$, has at most two types of parallel threads, entire and partial threads. Parallel threads of segment $S_{i,j}$ (other than the one in the master thread) have a maximum density $\delta_{i,j}^{max}$ of:

$$\delta_{i,j}^{max} = \max\left\{ \frac{e_{i,j}}{1 + f_{i,j} * (2 - 1)}, \frac{1 + (f_{i,j} - [f_{i,j}])e_{i,j}}{1 + f_{i,j}} \right\}$$

$$\forall m_{i,j} > 1 \Rightarrow \delta_{i,j}^{max} = \frac{1}{1 + f_{i,j}} \quad (7)$$

From the above equation, the maximum density of threads of segment $S_{i,j}$ in $\tau_i$ depends on the segment distribution factor $f_{i,j}$. From Equation 3, $f_{i,j}$ depends on the number of threads in this segment. Hence, the segment with a small number of threads has a higher density. We can conclude that the highest density of a thread in task $\tau_i$ occurs when the segment has the smallest possible number of threads ($m_{i,j} \geq 2$).

$$\delta_{i}^{max} = \frac{1}{1 + (f_i * (2 - 1))} \quad \rightarrow \text{ from 3}$$

$$= \frac{1}{1 + f_i} \quad (8)$$

From Equation 2, we can derive the following relation:

$$\forall f_i \geq 0 \rightarrow \delta_{i}^{max} = \frac{1}{1 + f_i} \quad \text{and} \quad f_i \geq \frac{D_i - L_i}{C_i}$$

$$\Rightarrow \frac{1}{1 + f_i} \leq \frac{1 + \frac{D_i - L_i}{C_i}}{1 + \frac{D_i - L_i}{C_i}} \leq \frac{C_i}{C_i + D_i - L_i} \leq 1$$

(Since $L_i \leq D_i$).

From both cases, the maximum density $\delta^{max}(\tau'')$ of all threads in $\tau''$ is:

$$\delta^{max}(\tau'') = \max_{\tau'' \in \tau''} \{\delta^{master}, \delta_i^{max}\}$$

$$\leq 1 \quad (9)$$

It is worth noticing here that DAG tasks with utilization equal to one are not considered in the previous cases, because the DAG stretching algorithm transforms them into fully-stretched master threads that are assigned their own executing processors. Hence, they are not included in the resource augmentation bound of global EDF scheduling.

In order to calculate the total density of thread set $\tau''$, we consider that at time $t$, only a single segment from each task with utilization higher than 1 can be active in addition to the master threads of tasks with utilization $< 1$. We calculate the density of a task $\tau_i''$, with $U_i > 1$, by considering the threads of the segment with the highest density. Let $\delta_i^{sum}$ be the sum of densities of such threads in a certain segment in a task $\tau_i$:

$$\delta_i^{um} \leq \frac{1}{1 + f_{i,j}} \ast (m_{i,j} - 1 - [f_{i,j}]) \quad \rightarrow \quad [f_{i,j}] \geq 0$$

and from Eq. (5)

$$\leq \frac{m_{i,j} - 1}{f_i \ast (m_{i,j} - 1)}$$

$$\leq \frac{1}{f_i}$$

$$\leq \frac{C_i - L_i}{D_i - L_i} \leq \frac{C_i}{D_i - L_i} \quad \rightarrow \text{ from Eq. (2)}$$

(10)

Now, consider the second case where task $\tau_i''$ has a utilization less than 1 and is transformed into an implicit-deadline non-fully stretched master thread. The total density of the task is denoted $\delta_{i}^{master} = \frac{C_i}{D_i} \leq \delta_{i}^{sum}$.

For taskset $\tau''$, let $\delta^{sum}(\tau'')$ be the sum of densities of every DAG task in the original taskset $\tau$:

$$\delta^{sum}(\tau'') \leq \sum_{\tau_i \in \tau} \frac{C_i}{D_i - L_i} \quad (11)$$

In order to calculate the speedup factor, we consider that taskset $\tau''$ executes on $m'$ processors with a minimum speed $\nu$, where $\nu > 1$. Increasing the speed of processors affects the execution parameters of task $\tau_i$ such as $C_i$ and $L_i$. The critical path length $L_i$ of any task $\tau_i$ has a necessary feasibility condition:

$$\forall \tau_i : \ 1 \leq i \leq n \quad \rightarrow \quad L_i \leq D_i$$

(12)

Otherwise, task $\tau_i$ is not feasible on a unit-speed processors.

On a processor that is $\nu$ times faster, the critical path length $L_i''$ is given by:

$$\forall \tau_i : \ 1 \leq i \leq n \quad \rightarrow \quad L_i'' = \frac{L_i}{\nu} \leq \frac{D_i}{\nu}$$

(13)

The same notation is applied for the total execution time $C_i''$ of task $\tau_i$ when it is run on processors $\nu$ times faster:

$$\forall \tau_i : \ 1 \leq i \leq n \quad \rightarrow \quad C_i'' = \frac{C_i}{\nu}$$

(14)

The total density of taskset $\delta^{sum,\nu}$ on speed-$\nu$ processors is:

$$\delta^{sum,\nu}(\tau'') \leq \sum_{\tau_i \in \tau} \frac{C_i''}{D_i - L_i''} \quad (15)$$

From Equation 13,

$$\forall 1 \leq i \leq n \quad \rightarrow \quad L_i'' \leq \frac{D_i}{\nu} \Rightarrow (D_i - L_i'') \geq D_i(1 - \frac{1}{\nu})$$

Using this result,

$$\delta^{sum,\nu}(\tau'') \leq \sum_{\tau_i \in \tau} \frac{C_i / \nu}{D_i(1 - \frac{1}{\nu})}$$

$$\leq \frac{1}{\nu - 1} \sum_{\tau_i \in \tau} \frac{C_i}{D_i}$$

$$\leq \frac{m}{\nu - 1} \quad (m)$$
From Equation 9, the maximum density $\delta^{\text{max},\nu}(\tau'')$ of taskset $\tau''$ on speed-$\nu$ processors is:

$$\delta^{\text{max},\nu}(\tau'') = \frac{\delta^{\text{max}}}{\nu} \leq \frac{1}{\nu} \quad (16)$$

Therefore, taskset $\tau''$ is schedulable under preemptive global EDF on $m'$ speed-$\nu$ processors if:

$$\frac{m}{\nu - 1} \leq \frac{m' + n}{\nu - 1} \leq \frac{m'}{\nu - 1} - (m' - 1) \frac{1}{\nu}$$

$$\frac{m'}{\nu - 1} + \frac{n}{\nu - 1} \leq \frac{m'}{\nu - 1} + \frac{1}{\nu} \leq 1 - \frac{n}{(\nu - 1)m'}$$

$$\frac{2\nu - 1}{\nu^2 - \nu} \leq 1 - \frac{n}{m'(\nu - 1)}$$

In order to have a positive value for the right hand side of the inequality, we consider the following:

$$0 < \frac{n}{m'(\nu - 1)} < 1 \Rightarrow m' > \frac{n}{\nu - 1}$$

$$\frac{2\nu - 1}{\nu^2 - \nu} < 1$$

$$0 < \nu^2 - 3\nu + 1$$

$$\Rightarrow \nu = \frac{3 + \sqrt{5}}{2} \quad \text{for} \ n < \nu \cdot m', \text{where} \ \nu = \frac{1 + \sqrt{5}}{2}$$

where $\varphi$ is the golden ratio.

This concludes the proof of the speedup factor. \square

6. SIMULATION

In this section, we provide simulation results for global EDF scheduling with stretched tasksets. The objective of simulation experiments is to test the performance of our DAG stretching algorithm when used prior to the scheduling process of DAG tasks. Also, we aim to compare it with another existing DAG scheduling technique found in the state-of-the-art [12]. To our knowledge, the decomposition algorithm proposed in [12] is the only other DAG scheduling technique found in the state-of-the-art. To our knowledge, the decomposition algorithm proposed in [12] is the only other DAG scheduling technique found in the state-of-the-art. In all the experiments, the percentage of schedulable tasksets when global EDF is used for $m$ processors (from 2 to 16). The x-axis of the figure represents the percentage of successfully schedulable tasksets when global EDF is used for $m$ processors (from 2 to 16).

![Figure 4: Comparison results of global EDF scheduling simulation between the DAG stretching algorithm and decomposition algorithm from [12]](image)

Comparison between DAG stretching and decomposition algorithms

The first results of simulation are provided in Figure 4. They are obtained by varying two parameters, the number of unit-speed processors $m$ on which tasksets are executing, and the percentage of taskset utilization, where the maximum utilization of a taskset is equal to $m$. Figure 4 shows the comparison results between our DAG stretching algorithm (denoted ‘STR’ in the figure) and the DAG decomposition algorithm $\delta^n$ (denoted by ‘DCMP’ in the figure). Similarly, The decomposition algorithm transforms each dependent DAG task into a set of independent sequential threads with intermediate offsets and deadlines by distributing the slack of the DAG on its segments. Figure 4 shows the percentage of successfully schedulable tasksets when global EDF is used for $m$ processors (from 2 to 16). The x-axis of the figure represents the percentage of taskset utilization w.r.t. the maximum utilization which is equal to $m$. The percentage values are between 20% and 100%. For each taskset, the scheduling process is simulated for both cases, when the taskset is stretched (STR) or decomposed (DCMP).

From our simulations, we notice that the performances of our stretching algorithm is better than the decomposition algorithm for the scheduling of the DAG tasks. In all the experiments, the percentage of schedulable tasksets when stretched are higher than the decomposed ones. Also, we can notice that the percentage of the schedulable tasksets decreases when the number of processors increases.

Varying the speed of processors

Here, we analyze the effect of increasing the speed of processors w.r.t. the schedulability of stretched DAG tasks. Based on the speedup factor of our stretching approach, the schedulability of stretched DAG tasks increases when the speed of processors increases and feasible tasksets become schedulable using global EDF when the speed $\nu > \frac{1 + \sqrt{5}}{2}$. In these experiments, we increased the speed of processors from $\nu = 1$ up to 4 by steps of 1, while varying the number of processor and taskset utilization as in the above simulations.

The results shown in Figure 5 are done by varying the

\footnote{For more details about the decomposition algorithm, please refer to the original paper in [12].}
speed of processors on which tasksets are executed. For each $m$ from 2 to 8, we simulate global EDF scheduling of tasksets of maximum utilization (equal to $m$). The choice of taskset utilization is done because these tasksets have the smallest success ratio in the results from Figure 4. As shown in Figure 5, schedulability of stretched tasksets increases when the speed of processors increases. In the case of $m = 2$ and $m = 4$, a success ratio of almost 100% is achieved when the speed of processors $\nu \geq 2$, while it needs processors speed at least equal to 4 for $m = 8$.

The simulation results provided in this section show the performances of our DAG stretching algorithm when compared to the decomposition scheduling algorithm of DAGs. Although simulation experiments based on random generation of tasksets cannot be used as an accurate performance metric of algorithms in real-time systems, the results are considered as an indication regarding the average behavior of studied scheduling algorithms.

7. CONCLUSIONS

In this paper, we presented a stretching algorithm for parallel periodic real-time DAG tasks on homogeneous multiprocessor systems. The stretching algorithm is an extension to the one in [9] which targeted Fork-join task model, a special case of DAG model. Similarly, our DAG stretching algorithm aimed at removing dependencies between subtasks of DAG tasks and stretching their critical path (master thread) up to its deadline. In addition to the master thread which is assigned its own executing processor, a set of independent threads with intermediate offsets and deadlines are generated and scheduled using global EDF scheduling algorithm.

Then, we proved a speedup factor of $\frac{\phi \sqrt{\phi}}{m'}$ for global EDF scheduling algorithm for stretched threads if $n < \phi \cdot m'$, where $n$ is the total number of tasks in the taskset, $m'$ is the remaining number of processors after taking off the processors dedicated for the fully-stretched master threads and $\phi$ is the golden ratio whose value is $\frac{1+\sqrt{5}}{2}$ and 4 else.

Finally, we compared the performance of our stretching algorithm with another scheduling technique for DAGs from the state-of-art using DAG decomposition. The simulation results show that our DAG stretching algorithm has better average performances.

8. REFERENCES